

A VAM AND MODI METHOD TO SOLVE THE OPTIMALITY FOR THE TRANSPORTATION PROBLEMSonia Shivhare*¹^{*1}Amity university Madhya Pradesh
sshivhare@live.com**Abstract:**

In this paper, an optimal solution of Transportation programming problem has been considered. To obtain initial basic feasible solution (IBFS), the work shown here is done with Vogel's approximation method (VAM). The MODI method is applied for finding an optimal solution for this T.P.P. The proposed method is unique and the result with an elaborate illustration demonstrates that the method presented here is effective optimal test of the transportation cost.

Keywords:

Transportation problem (T.P.), VAM, MODI Method, Initial basic feasible solution (IBFS), optimal test

INTRODUCTION

The Transportation problem is one of the traditional function of the L.P.P. Transportation model provides a greater impact on the management of transport [5]. The basic Transportation problem was initially proposed by Hitchcock [3][6]. It is a special kind of the network optimization problems. It has defined data structure in solution, characterized as a Transportation group[1]. It plays an important role in logistics & supply chain management for reducing cost & improving service[4]. The problem mainly deals with the determination of optimum routes to minimize transporting cost of particular product from a number of sources to a number of destinations. Network model of the transportation programming problem is shown in figure 1. It aims to find the best method to fulfill the demand of n-demand points using the capacities of m-supply points[7] by applying the renowned method i.e., Vogel's Approximation Method (VAM) and check its optimality by applying Modified Distribution (MODI) test method.

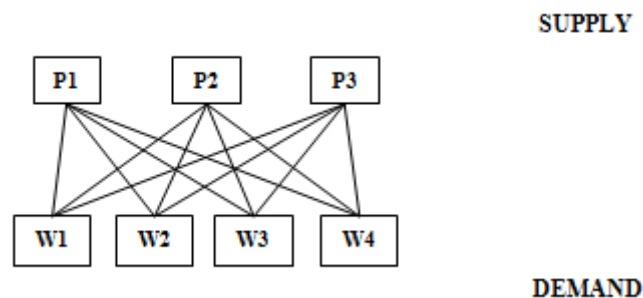


Figure: 1 - Network Model of T.P.P.

ALGORITHM OF VOGEL'S APPROXIMATION METHOD

- Step 1:** Verify whether the total supply equals total demand. If it is not so then set up a dummy row or column.
- Step 2:** Identify least and second least transportation cost in each row or column and write the difference (penalty) beside the corresponding row or column.
- Step 3:** Now select the row or column with maximum penalty and make maximum allotment to the box having minimum cost of transportation in that row or column. If the penalties of two or more rows or columns are equal, then choose any one of them.
- Step 4:** Adjust the supply and demand and eliminate the satisfied row or column. If a row and column are satisfied simultaneously, only one of them is eliminated and the other one is assigned a nil value. Any row or column having nil supply or demand cannot be used in calculating future penalties. [8]
- Step 5:** Repeat the process until all the supply sources and demand destinations are satisfied.

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ILLUSTRATION

To / From	D ₁	D ₂	D ₃	SUPPLY
S ₁	2	7	4	5
S ₂	3	3	1	8
S ₃	5	4	7	7
S ₄	1	6	2	14
DEMAND	7	9	18	34 / 34

Table 1

Solution: By Using VAM Method

To / From	D ₁	D ₂	D ₃	SUPPLY	PENALTY (P ₁)	P2	P3	P4	P5
S ₁	2 5	7	4	5	2
S ₂	3	3	1 8	8	2	2
S ₃	5	4 7	7	7	1	1	1	1	4
S ₄	1 2	6 2	10	14 4 2	1	1	1	5	6
DEMAND	7 2	9 7	18 10	34 34					
PENALTY (P ₁)	1	1	1						
P2	2	1	1						

P3	4	2	5
P4	4	2
P5	2

Table 2

From this table, it can be seen that the number of non-negative independent allocations is

$$(m+n-1) = (4+3 - 1) = 6. [5]$$

Hence the solution is non-degenerate basic feasible.

The initial transportation cost: $5*2 + 2*1 + 7*4 + 2*6 + 8*1 + 10*2 = \text{Rs.}80$

TEST THE SOLUTION FOR OPTIMALITY

By Modified Distribution Method (MODI Method) or u-v Method:

To verify whether the feasible solution obtained minimizes the total transportation cost or not .To find optimal solution of transportation problem we apply MODI (Modified Distribution method) test as follows:

Step-1: Determine an initial basic feasible solution as we derived in the above problem by Vogel Approximation Method. Then take the costs only of the occupied cells.

Step-2: Set $v_1, v_2, v_3 \dots$ etc. against the corresponding column and set $u_1, u_2, u_3 \dots$ etc. against the corresponding row. Then determine a set of u_i and v_j s.t. for each occupied cell, using $u_i + v_j = c_{ij}$. [2]

Step-3: Assign 0 to one of the u_i 's or v_j 's for which the corresponding row or column have the maximum number of individual allocations.

Step-4: Then find the cell evaluations (Subtract the above matrix's cells from the corresponding cells of original matrix) $u_i + v_j$ for each unoccupied cell and enter at the upper right corner of the corresponding unoccupied cell.

Step-5: Then we calculate d_{ij} (difference of each occupied cell) by the $d_{ij} = C_{ij} - (u_i + v_j)$ for each unoccupied cell and enter at the lower right corner of the corresponding unoccupied cell.

Step-6: If all the d_{ij} is non-negative, then the basic feasible solution is optimal. On the other hand, if anyone of d_{ij} is $-ve$, then the basic feasible solution is not optimal.

Step-7: Select the largest negative value of d_{ij} . If there have more than one equal cell, then any one can be chosen. Then draw a closed loop for the unoccupied cell. Starting the closed loop with the largest negative value of d_{ij} and draw a closed loop with the occupied cells only.

Step-8: Mark the identified cell as $+ve$ and each occupied cell at the corners of the path alternatively $-ve, +ve, -ve$ and so on.

Step-9: check all the negative position and consider the smallest transportation cost that has been assigned a $-ve$ sign. Now, $+$ and $-$ demand and supply values of all the positions of $+$ and $-$.

Step-10: Repeat the whole procedure until the optimum solution is obtained.

Solution:

To From	D ₁	D ₂	D ₃	a _j
S ₁	2 5	7	4	5
S ₂	3	3	1 8	8
S ₃	5	4 7	7	7
S ₄	1 2	6 2	2 10	14
b _j	7	9	18	

Table 3

Now we determine a set of u_i and v_j s.t. for each occupied cell, $c_{ij} = u_i + v_j$

For this we choose $u_4 = 0$ (since row 4 contains maximum no. of allocations).

Since $C_{41} = u_4 + v_1 = 1$ therefore, $v_1 = C_{41} - u_4 = 1$

$$C_{42} = u_4 + v_2 = 6 \quad v_2 = C_{42} - u_4 = 6$$

$$C_{43} = u_4 + v_3 = 2 \quad v_3 = C_{43} - u_4 = 2$$

$$C_{11} = u_1 + v_1 = 2 \quad u_1 = C_{11} - v_1 = 1$$

$$C_{23} = u_2 + v_3 = 1 \quad u_2 = C_{23} - v_3 = -1$$

$$C_{32} = u_3 + v_2 = 4 \quad u_3 = C_{32} - v_2 = -2$$

Then we find the cell evaluations $u_i + v_j$ for each unoccupied cell by $d_{ij} = c_{ij} - (u_i + v_j)$

(2) 5	(7) (7) (0)	(4) (3) (1)	1
(3) (0) (3)	(3) (5) (-2)*	(1) 8	-1
(5) (-1) (6)	(4) 7	(7) (0) (7)	-2
(1) 2	(6) 2	(2) 10	0

- Since cell evaluation $d_{22} = -2 < 0$, so the solution under test is not optimal.
- Since d_{22} has the largest negative value, so making a closed path with the help of this.
- Starting the closed loop with the largest negative value of d_{ij} and making a closed loop with the occupied cell only.

5		
	$+\Theta$	$8 - \Theta$
	7	
2	$2 - \Theta$	$10 + \Theta$

- Since minimum allocation containing $-\Theta$ is $2 - \Theta$
Taking $2 - \Theta = 0$, we get $\Theta = 2$
- The new B.F.S. thus obtained is shown in the following table :

(2) 5	(7)	(4)	5
(5)	(3) 2	(1) 6	8
(5)	(4) 7	(7)	7
(1) 2	(6)	(2) 12	14

Table 6

- Now the new transportation cost = $5*2+2*1+2*3+7*4+6*1+12*2 = \text{Rs. } 76$
- Now set u_1, u_2, u_3, u_4 and v_1, v_2, v_3 again.

u_i

(2) 5	(7) (5) (2)	(4) (3) (1)	1
(3) (0) (3)	(3) 2	(1) 6	-1
(5) (1)	(4) 7	(7) (2)	0
(4) (1) 2	(6) (4) (2)	(2) 12	0

Table 7

- Since all $d_{ij} > 0$. Hence, the B.F.S. shown by table is an optimum solution which is also unique.

CONCLUSION

The approach of VAM is simple in calculation and easy to understand and apply. The method discussed gives us an initial basic feasible solution of a balanced transportation problem in minimization cost. The method developed here gives the optimal solution.

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